

NIP formulas in continuous logic

Setting: M metric structure (in particular a complete bdd metric space), $G = \text{Aut}(M)$.

* If M is separable, it is called ω -categorical if unique separable model of $\text{Th}(M)$ up to isomorphism.

Recall (continuous Ayala - Nardeschi):

M separable. TFAE

- M is ω -cat.
- M^ω / G is compact ($G \sim \mathbb{N}$ approximately diffeomorphic)
- All uniformly continuous G -invariant bdd functions $h: M^\omega \rightarrow \mathbb{R}$ are given by formulas.

Terminology: M \emptyset -saturated if M realizes all $p \in S_\omega(\emptyset)$.

[Fact 1: M ω -cat $\Rightarrow M$ \emptyset -saturated.

Fix countable ordinals α, β , variables $(x_i)_{i < \alpha} = x$ and $y = (y_i)_{i < \beta}$ and a formula $\varphi(x; y): M^\alpha \times M^\beta \rightarrow \mathbb{R}$.

Moreover, let $A \subseteq M^\alpha$ and $B \subseteq M^\beta$ be \emptyset -type/definable sets. If $\tilde{M} \preceq M$, we write $\tilde{\varphi}$ and \tilde{A}, \tilde{B} for the interpretations in \tilde{M} .

Prop. 2 (Char of NIP formulas)

Let M be \emptyset -saturated and f, A, B as above. TFAE

(1) There are no $r \neq s$ in \mathbb{R} and $(a_i)_{i < \omega}$ from A and $(b_i)_{i < \omega}$ from \tilde{B} s.t. $\forall I \subseteq \omega$:

$$\lim_{i \in I} f(a_i, b_i) = \begin{cases} r & \text{if } i \in I \\ s & \text{if } i \notin I \end{cases}$$

[\cong no infinite set is shattered, in classical discrete logic]

(2) For any indiscernible $(a_i)_{i < \omega}$ in A and $b \in \tilde{B}$

$\lim_{i \rightarrow \infty} f(a_i, b)$ exists. [\cong finite alternation]

(3) For any sequence $(a_i)_{i < \omega}$ in A there is a subsequence

$(a_{i_j})_{j < \omega}$ s.t. $\lim_{j \rightarrow \infty} f(a_{i_j}, b)$ exists for all $b \in \tilde{B}$.

$f(x, y)$ is said to be NIP on $A \times B$ if these equivalent conditions hold.

Pf: (3) \Rightarrow (1): Trivial, since any subsequence $(a_{i_j})_{j < \omega}$ of

a counter to (1) yields $r = \lim_{j \rightarrow \infty} f(a_{i_j}, b_j) = s$

wherever $\mathcal{J} \subseteq \omega$ intersects $\{i_j \mid j \in \omega\}$ in an infinite coinitial set.

(2) \Rightarrow (3): By compactness, for every $\varepsilon > 0$ there is a finite set Δ of formulas and $\delta > 0$ s.t. whenever $(a_i)_{i \in \omega}$ is Δ - δ -indiscernible in A (i.e. $|g(a_{i_1}, \dots, a_{i_n}) - g(a_{j_1}, \dots, a_{j_n})| \leq \delta$ for any $g \in \Delta$ and $i_1 < \dots, i_n$ and $j_1 < \dots, j_n$ from ω) and $b \in \tilde{B}$, there is $N \in \omega$ s.t.

$$|\tilde{f}(a_i, b) - \tilde{f}(a_j, b)| < \varepsilon \text{ for all } i, j \gg n.$$

Now let (Δ_n, δ_n) correspond to $\varepsilon_n = \frac{1}{n}$.

③

- Set $(a_i^0)_{i < \omega} := (a_i)_{i < \omega}$.

- Given $n \in \omega$ and assuming $(a_i^n)_{i < \omega}$ is a subsequence of $(a_i)_{i < \omega}$, by Ramsey we find a $\Delta_{n+1}, \delta_{n+1}$ -indiscernible subsequence $(a_i^{n+1})_{i < \omega}$ of $(a_i^n)_{i < \omega}$.

Diagonalize! Set $a_{i_j} := a_i^j$.

By construction $(a_{i_j})_{j < \omega}$ is as we wanted, i.e. $\lim_{j \rightarrow \infty} f(a_{i_j}, b)$

exists for all $b \in \tilde{B}$.

① \Rightarrow ②: If $f(a_i, b)$ does not converge, for $(a_i)_{i < \omega}$ indiscernible in A at $b \in \tilde{B}$, we find $i_0 < j_0 < i_1 < j_1 < \dots$ in ω and $r \neq s$ in \mathbb{R} s.t. $\tilde{f}(a_{i_k}, b) \xrightarrow{k \rightarrow \infty} r$ and $\tilde{f}(a_{j_k}, b) \xrightarrow{k \rightarrow \infty} s$.

~~Passing~~ Pass to a subsequence, w.o.a. $\tilde{f}(a_{2k}, b) \rightarrow r$ and $\tilde{f}(a_{2k+1}, b) \rightarrow s$.

\mathcal{M} ϕ -saturated \Rightarrow w.o.a. $\tilde{f}(a_{2k}, b) = r$ and $\tilde{f}(a_{2k+1}, b) = s$ $\forall k$.

Now let $I \subseteq \omega$ be arbitrary. Let $\tau: \omega \rightarrow \omega, \tau(n) := \begin{cases} 2n & \text{if } n \in I \\ 2n+1 & \text{else} \end{cases}$

$\sigma(a_i) := a_{\tau(i)}$ defines a partial elementary map, extend it to

$\tilde{\sigma} \in \text{Aut}(\tilde{\mathcal{M}})$. Let $b_{\pm} := \tilde{\sigma}^{-1}(b)$. Then $\tilde{f}(a_i, b_{\pm}) = \begin{cases} r & \text{if } i \in I \\ s & \text{if } i \notin I \end{cases}$ by construction. \blacksquare

RK3: ③ in Prop. 2 is equivalent to the following:

Let $B^* := \{ \kappa(b/M) \in S_y(M) \mid b \in \tilde{M} \subset M \}$. Then $\{ \tilde{f}_a|_{B^*} : a \in A \}$ is a sequentially precompact subset of $\mathcal{C}(B^*)$. Then

$\{ \tilde{f}_a|_{B^*} : a \in A \}$ is a sequentially precompact subset of $\mathcal{C}(B^*)$, where $\tilde{f}_a|_{B^*}$ is the function on B^* induced by $\tilde{f}(q, y)$.

[I.e., every sequence admits a convergent subsequence.]

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Link to tame functions

Let X be a G -space

$\mathcal{C}(X) := \{ f: X \rightarrow \mathbb{R} \text{ continuous + bdd} \}$

\exists Defⁿ of a "tame" function $f \in \mathcal{C}(X)$ (we don't give it)

[Fact: If X is compact, $f \in \mathcal{C}(X)$ is tame iff $G \cdot f$ is sequentially precompact in \mathbb{R}^X .

Here, $(g \cdot f)(x) := f(g^{-1} \cdot x)$

[Prop 5: M w. act, $f(x, y)$ tame, $B \subseteq M^B$ (M -) definable, $a \in M^a$, then $\tilde{f}_a \in \mathcal{C}(B^*)$ is tame iff $f(x, y)$ is NIP on $[a] \times B$.

Note that $[a]$ is closed and G -invariant, so (M) definable. As $G \cdot a$ is dense in $[a]$, $\{ \tilde{f}_{a'}|_{B^*} : a' \in [a] \}$ is sequentially precompact iff

$\{ \tilde{f}_{g \cdot a}|_{B^*} : g \in G \}$ is, so we conclude by Remark 3 + Prop 2.

Now let $B = M$ (or $B = M^\beta$ for some c.t.b. β).

The previous proposition, characterizing tame functions, applies to the set of fns

(5)

$$\text{Def}(M) := \{f(a, y) : M \rightarrow \mathbb{R} \mid f(x, y) \text{ formula, } a \in M^{\omega}\} \subseteq \mathcal{C}(X)$$

Given a G -space X , we set

$$\text{RUC}(X) := \{f \in \mathcal{C}(X) \mid \forall \epsilon > 0 \exists U \text{ nbhd of } 1 \in G : \|g \cdot f - f\| < \epsilon \forall g \in U\}$$

If X is a metric space,

$$\text{RUC}_u(X) := \text{RUC}(X) \cap \{f : X \rightarrow \mathbb{R} \mid f \text{ unif cont}\}$$

Lemma 6: M ω -cat metric structure, $G = \text{Aut}(M)$. Then

$$\text{Def}(M) = \text{RUC}_u(M)$$

Pf sketch: " \subseteq " holds in any metric structure (easy exercise)
ad \supseteq : Let $h \in \text{RUC}_u(M)$. Let $a \in M^\omega$ enumerate a dense subset of M . Define $f : G \cdot a \times M \rightarrow \mathbb{R}$ via

$$f(ga, b) := (g \cdot h)(b) = h(g^{-1}b)$$

Since a is dense in M , f is well-def. Moreover, f is bdd, unif cont + G -invariant (for diagonal G -action on $G \cdot a \times M$).

[Exercise: use $h \in \text{RUC}_u(M)$ + a dense in M]

(i) f extends to $\tilde{f} : [a] \times M \rightarrow \mathbb{R}$ unif cont bdd G -invar

(ii) \tilde{f} factors throug $\bar{f} : ([a] \times M) // G \rightarrow \mathbb{R}$ unif cont + bdd

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 (iii) \bar{f} extends to $\bar{F} : (M^\omega \times M) // G \rightarrow \mathbb{R}$ unif cont + bdd.

(iv) $\bar{F} = \bar{F} \circ \pi : M^\omega \times M \rightarrow \mathbb{R}$ unif cont + bdd + G -invar

Since π is ω -cat, $\bar{F}(x, y)$ is a formula, and by construction

$$h(y) = \bar{F}(a, y), \text{ so } h \in \text{Def}(M)$$

