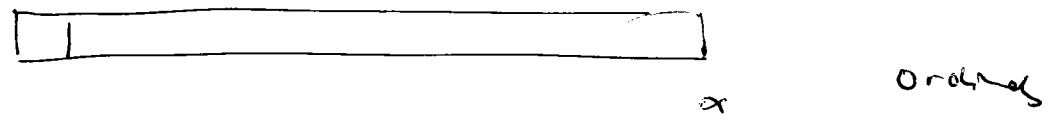


Philipp Schlicht,

recognizable sets of ordinals.

(joint with merlin c2l)

ordinal Turing machines OTM



lim tree  $\lambda$  lim state

- ①  $p$  states  $\mapsto$  states
- ② content of the cells.
- ③ head position: lim inf

$x \subset \alpha$  (OTM-) Computable for  $\alpha$

$(\Rightarrow)$   $x$  is an output of a computation which has  $\alpha$  as an additional input.

$x$  computable  $\iff x \in L$ .

$x$  is recognizable iff there is a program  $P$  and some  $\beta$

$$\forall y \in \alpha$$

$$(P(y, \beta) = 1 \iff x = y)$$

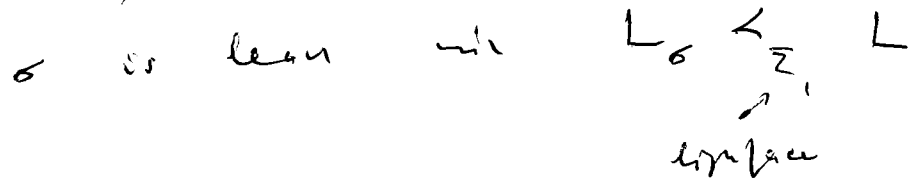
?  
output of computer  
with input  $y$  and  
 $\beta$  as add. input.

without additional ordinal parameter:

$x \in \omega$  computable iff

$x$  is recognizable  $\iff$

$x \in L_\sigma$ , where



assume that  $M_1^\#$  exists.

question. is every recognizable in  
of ordinals in  $L$ ?

def.  $x \subset \alpha$  is in the  
 recognizable closure if there are  
 $x = x_0, x_1, \dots, x_n$ ,  $x_n$  recognizable,  
 $x_i$  recognizable for  $x_{i+1}$ .

lem. Suppose that  $x \subset \alpha$ . TFAE.

$x$  is constructible from  $y \subset \beta$

1.  $y$  is recognizable

2. there is for  $\varphi$ , some  $\beta \gamma$  s.t.

$y$  is the unique subset  $Z$  of

$\beta$  s.t.  $L[Z] \models \varphi(Z, \gamma)$ .

3.  $y$  implicitly defines over  $L$ ,

i.e. there is for  $\varphi$  etc. and some  $\gamma$  s.t.

$y$  is the unique  $z \subset \beta$  s.t.

$(L, \epsilon, z) \models \varphi(z, \gamma)$

$M^\alpha =$  the  $\alpha$ <sup>th</sup> iterate of  $M$ ,  
 by the least number cardinal  
 + its image

$$\text{lem } M^\infty = \bigcup_{\alpha \in \text{OR}} M^\alpha$$

lem. if  $x \subset \alpha$  is recognizable, then  $x \in M^\infty$ , or  $\alpha < \omega_1$ .

question. is every recognizable subset  $\omega_1$  in  $M_1$ ?

lem. if TP is homogeneous,  $G$  is TP-generic over  $V$ .

supp. that  $y \subset \text{OR}$  is recognizable,  $y \in \text{VEGS}$ . then  $y \in V$ .

lem. supp.  $M_1^\#$  ex.,  $\text{ZF} + \text{DC}$  holds,

any  $X \subset \omega_1$  is in  $L[y]$ , for  $y \in \omega$ .

let TP be homogeneous, TP doesn't con.  $\omega_1$ , TP forces choice.

then any recognizable subset of  $\omega_1$  is in  $M_1$ .

question. supp. that  $H_{w_2}$  is

closed  $\rightarrow$   $M_1^\#$  ?

then is any rec. subset of  
of  $w_1$  in  $M_1$  ?

question. can we add recognizable

sets by forcing any large cardinals ?

lem. (steel) supp.  $\alpha < w_1^M$ .

say  $p_w(N) = w$ ,  $N$  has height  
 $\alpha$ ,  $N$  sound,

$N$  is  $\pi_2^1$  sl. th. then  $N = M_1 \parallel \alpha$ .

lem. supp. that  $\alpha < w_1^{M_1}$ ,

$p_w(M_1 \parallel \alpha) = w$ .

let  $y \in w$  be a can. code for

$M_1 \parallel \alpha$ . then  $y$  is recognizable.

then if  $x \in M_1 \parallel \alpha$ ,

$x$  is constructible from  $y$ .

$\Rightarrow M_1 \parallel w_1^M$  is recognizable in  $w_1^M$ .